Today: Single-source shortest paths. 24. $\{0,1,2,3\}$.
Next class: Homework $\mathbf{4}$ due. All-pairs shortest paths. $25 .\{0,1,2\}$.
Reminders: Read material before and after class. Use the class newsgroup. Quiz soon.

1. List the members of your group below. Underline your name.
2. Quick self-check. (Use textbook conventions.)
(a) If there is no path from $A$ to $B$, the shortest-path weight $\delta(A, B)$ is [circle the correct answer]: (1) 0 ; (2) -1 ; (3) $\infty$; (4) undefined.
(b) In order to use a single-source shortest paths algorithm to solve a single-destination shortest paths problem, we must: (1) negate edge weights (only); (2) reverse edges (only); (3) both negate edge weights and reverse edges; (4) use a completely different algorithm.
(c) True or false: If $\langle A, B, C, D, E\rangle$ is a shortest path from $A$ to $E$ then running Dijkstra's shortest path algorithm with source vertex $B$ must produce $\langle B, C, D\rangle$ as a shortest path from $B$ to $D$, due to the optimal substructure of the problem.
(d) Annotate each vertex in the following graph with the shortest-path weight from $A$ (or indicate that it is undefined).

(e) A shortest-paths algorithm can limit attention to cycle-free paths only if [circle all valid cases]: (1) there are no negative-weight edges; (2) there are no negative-edge cycles; (3) there are no cycles; (4) there are no unreachable vertices.
(f) If a a graph has negative-weight edges but no negative-weight cycles, shortest paths may be computed using [circle all correct options]: (1) Dijkstra's algorithm; (2) Bellman-Ford algorithm; (3) Depth-first search; (4) none of the above.
(g) The predecessor subgraph induced by the $v . \pi$ values produced by algorithms in this chapter (always): [circle all correct options]: (1) is connected; (2) is a tree; (3) is a DAG; (4) none of the above.
3. Trace the execution of the Bellman-Ford single-source shortest paths (SSSP) algorithm on the following directed graph, with vertex A as the source. Use the textbook's Fig. 24.4 (p. 652) as a model. Relax edges in lexicographic order. Annotate predecessor edges with check marks.

[additional space for answering the earlier question]
4. Repeat Question 3 using Dijkstra's SSSP algorithm, treating all negative edge weights as positive, and using the textbook's Fig. 24.6 (p.659) as a guide.
[additional space for answering the earlier question]
5. Informal homework: Augment the answer to Question 4 with the states of the primary data structures at each step (pairing heap, union-find).
